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Published in:
SIAM Journal on Matrix Analysis and Applications

DOI:
[10.1137/090780493](https://doi.org/10.1137/090780493)

Publication date:
2011

Document Version
Early version, also known as pre-print

[Link to publication](#)

Citation for pulished version (HARVARD):
Gratton, S, Toint, P & Tshimanga Ilunga, J 2011, 'Range-space variants and inexact matrix-vector products in Krylov solvers for linear systems arising from inverse problems', *SIAM Journal on Matrix Analysis and Applications*, vol. 32, no. 3, pp. 969-986. <https://doi.org/10.1137/090780493>

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Range-space variants and inexact matrix-vector products in Krylov solvers for linear systems arising from inverse problems

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18 March 2011

Abstract

The object of this paper is to introduce range-space variants of standard Krylov iterative solvers for unsymmetric and symmetric linear systems, and to discuss how inexact matrix-vector products may be used in this context. The new range-space variants are characterized by possibly much lower storage and computational costs than their full-space counterparts, which is crucial in data assimilation applications and other inverse problems. However, this gain is achieved without sacrificing the inherent monotonicity properties of the original algorithms, which are of paramount importance in data assimilation applications. The use of inexact matrix-vector products is shown to further reduce computational cost in a controlled manner. Formal error bounds are derived on the size of the residuals obtained under two different accuracy models, and it is shown why a model controlling forward error on the product result is often preferable to one controlling backward error on the operator. Simple numerical examples finally illustrate the developed concepts and methods.

Keywords: Krylov methods, linear systems, inexact matrix-vector products, data assimilation.

1 Introduction and motivation

Inverse problems in the natural sciences and elsewhere often give rise to very large under-determined parameter identification problems. In these problems, one typically tries to identify system parameters by fitting a model's output to a number of observations which is very often much smaller than that of the parameters. The resulting under-determined fitting problem is then regularized by selecting parameter sets that are as close as possible to values known from a previous case study (see Gratton, Lawless and Nichols, 2007). As we discuss below, problems of this type often lead, possibly after preconditioning, to solving variational formulations for which one considers the iterative solution of (potentially very large) linear systems of the form

$$(\gamma I_n + K^T L)s = b, \quad (1.1)$$

where $\gamma I_n + K^T L$ is a nonsingular $n \times n$ matrix, both K and L are $m \times n$ matrices (possibly identical) with $m \ll n$, I_n is the identity matrix of size n and b is a general right-hand side. The objective of this paper is to propose and analyze iterative methods whose storage requirements and linear algebra costs (beyond that of the products of vectors by K^T and

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L) are essentially dependent on m (at variance with standard approaches where they vary with n), and which produce monotonically decreasing values of the underlying variational objective. Moreover, their computational cost is further reduced by considering that the product of a vector by K^T or L may be only approximate (see Simoncini and Szyld, 2003, van den Eshof and Sleijpen, 2004, van den Eshof, Sleijpen and van Gijzen, 2005, Giraud, Gratton and Langou, 2006).

The motivation for investigating problem (1.1) and methods of the type we just mentioned finds its origin in (but is not limited by) the data assimilation problem for oceanography and weather forecasting (see Rabier, 2005, for instance). In this important application, which we now describe, the model is that of a discretized Navier-Stokes formulation (or a variant thereof) for describing the ocean or atmosphere and the observations are given by temperature, current, pressure and altimetry measurements collected from satellites, ocean-bound or ground stations. The most commonly studied problem is then to determine a complete set of initial conditions from which the Navier-Stokes model can be integrated to provide predictions, with the difficulty that they are many fewer observations (despite today's large data collection capabilities) than variables to determine in this initial condition, thereby suggesting regularization techniques. In daily practice (2009), the number of such variables is of the order of 10^8 while the number of available observations is only of the order of 10^5 (see Rabier, 2005 again), and the ratio of the first to the second is furthermore expected to increase significantly with the advent of more complex procedures (see Trémolet, 2006a, 2006b). The previously known state from which deviation is minimized (for the purpose of regularization) is called the background state and often selected as the values of the unknown parameters at the current time derived from a model validated in the past. As explained in Gratton et al. (2007), one is then interested in solving the nonlinear least-squares problem given by

$$\min_x \frac{1}{2}(x - x_b)^T B^{-1}(x - x_b) + \frac{1}{2} \sum_{j=0}^p \{ \mathcal{H}_j[x(t_j)] - y_j^o \}^T R_j^{-1} \{ \mathcal{H}_j[x(t_j)] - y_j^o \}, \quad (1.2)$$

where x_b is the background state, $x = x(t_0)$ is the initial state of the system, \mathcal{H}_j is the operator modelling the observed quantities at time t_j , y_j^o is the vector of observations, and $x(t_j)$ the state at this time, and the matrices B and R represent correlations between the background variables and the observations, respectively. Problem (1.2) is typically solved by a standard truncated Gauss-Newton method (known in the oceanography and weather-forecasting communities under the name of incremental 4D-Var (Courtier, Thépaut and Hollingsworth, 1994), an acronym for four dimensional variational assimilation). The linearized subproblem arising at iteration k of this Gauss-Newton procedure, where we have concatenated the observations and model predictions over time into a single vector (see Gratton and Tshimanga, 2009 for details) is then given by

$$\min_s \frac{1}{2}(x_k + s - x_b)^T B^{-1}(x_k + s - x_b) + \frac{1}{2}(\mathcal{H}s - d)^T (\mathcal{H}s - d), \quad (1.3)$$

where x_k and x_b are known, x_k , x_b and s belong to \mathbb{R}^n with n being the dimension of the unknown initial state, B is a positive-definite $n \times n$ symmetric matrix, $d \in \mathbb{R}^m$ is the concatenated misfit vector at x_k multiplied by the square root of the inverse of R , where $R = \text{diag}(R_0, \dots, R_p)$ is $m \times m$ and symmetric positive-definite⁽¹⁾, and \mathcal{H} is a $m \times n$ matrix representing the concatenated linearized model also multiplied by the square root of the inverse of R . Note that this standard formulation does not introduce vectors or matrices of size $m + n$, a very desirable property given the magnitude of n . The minimization of the convex quadratic given by (1.3) would seem a reasonably well-mastered problem, if it were not for the sizes involved (recall that we aim at $n \approx 10^9$ and $m \approx 10^5$). As for smaller problems, today's standard method for the solution of (1.3) is to apply a

⁽¹⁾A more elaborate way to efficiently handle the correlation matrix R is of course used in practice, but we choose this implicit formulation here for the sake of simplicity.

conjugate-gradient algorithm, which may be viewed as an iterative solver for the system

$$(B^{-1} + \mathcal{H}^T \mathcal{H})s = \mathcal{H}^T d + B^{-1}(x_b - x_k). \quad (1.4)$$

As is typical for conjugate-gradient method, preconditioning is crucial for efficient numerical performance, and we consider here two distinct possibilities. The first preconditioning technique, which is closer to current practice employs $B^{1/2}$. It is two-sided and symmetric, and gives the system

$$(I_n + B^{1/2} \mathcal{H}^T \mathcal{H} B^{1/2})z = B^{1/2} \mathcal{H}^T d + B^{-1/2}(x_b - x_k) \text{ where } s = B^{1/2} z. \quad (1.5)$$

Interestingly, the cost of the products with \mathcal{H} , which involves the solution of the time dependent partial-differential equation defining $x(t_j)$ in (1.2), is so high in practical applications that every effort is made to reduce the number of such products beyond the effect of preconditioning. This even includes the use of reorthogonalization within the considered iterative solvers, which often implies large storage requirements. As a result, the interest of maintaining a purely symmetric formulation becomes less clear. This makes a second preconditioning technique attractive, where one uses a simple right-sided product by B , yielding the system

$$(I_n + \mathcal{H}^T \mathcal{H} B)z = \mathcal{H}^T d + B^{-1}(x_b - x_k) \text{ where } s = Bz. \quad (1.6)$$

It is also useful to note that, in the framework of the original nonlinear problem (1.2), the vector $x_b - x_k$ can often be itself written in the form $B\mathcal{H}^T u$ for some vector u in the range of \mathcal{H} given by the step from the background state x_b to the current nonlinear iterate (see Gratton and Tshimanga, 2009). This in turn implies that $B^{-1}(x_b - x_k) = \mathcal{H}^T u$ and the complete right-hand side of (1.5) then lies in the range space of $K^T = \mathcal{H}^T$.

Our motivation may now be understood since both (1.6) and (1.5) are *examples*⁽²⁾ of systems of the form (1.1), and that their practical solution by efficient iterative methods correspond to the declared objective of this paper. It should nevertheless be stressed that the numerically efficient algorithms for the solution of (1.1) are of interest much more broadly, as this framework covers a large class of (mostly inverse) problems considered in a variational setting and where the “action” is limited to a small subspace of the original formulation. The standard Tikhonov regularization scheme for rank-deficient or ill-posed problems also leads to systems of the form (1.1) (see Hansen, 1997), in which case $\gamma > 0$ is typically different from 1. A more complete discussion of the effect of using iterative methods in the context of data assimilation is also of interest, where one has to consider the possible but hopefully limited contamination of the solution with noise arising from the approximate solution of the involved systems. This analysis may fortunately be conducted and shows that the use of iterative method can be both accurate and efficient in this specific context. This non-trivial analysis is however outside the scope of the present paper and will be reported on elsewhere.

As can be seen from the minimization context introduced, it is also crucially desirable for the design of realistic termination criteria that iterates produced by the solution algorithms ensure an *iteration-wise decrease of the underlying objective function* in (1.3), in the symmetric case, or *of the associated residual norm* otherwise. Unfortunately, a simple solution consisting in using the Sherman-Morrison-Woodbury formula⁽³⁾ on the linear system and then applying an Krylov-space iterative method on the resulting $m \times m$ linear system, or variants thereof such as the PSAS method⁽⁴⁾ or the equivalent “representer” technique of Bennett and Thornburn (1992), may often be problematic. The main difficulty⁽⁵⁾ is that the underlying objective function (for the symmetric case) or the residual

⁽²⁾In (1.6), we have that $K = \mathcal{H}$ and $L = \mathcal{H}B$, while $K = L = \mathcal{H}B^{1/2}$ in (1.5).

⁽³⁾See Conn, Gould and Toint, 2000, page 57, for instance.

⁽⁴⁾See Amodei (1995), Da Silva, Pfaendtner, Guo, Sienkiewicz and Cohn, 1995, and Cohn, Da Silva, Guo, Sienkiewicz and Lamich, 1998.

⁽⁵⁾See El Akkrouri, Gauthier, Pellerin and Buis, 2008, or Gratton and Tshimanga, 2009.

norm (for the unsymmetric case) are not monotonic in the original variables for these approaches, and therefore cannot be used in convenient and robust termination criteria. As a result, the iterations must often be carried out much longer than necessary.

Our developments originate in the recent proposal by Gratton and Tshimanga (2009), where a variant of the conjugate-gradient method, called RPCG, was proposed for the exact symmetric case, i.e. the case where $L = K$ and the products by K are carried out exactly. As for the variants discussed below, RPCG's storage requirement and linear algebra costs (beyond that of the products by K and K^T) vary with m rather than n . However, and although initial numerical experiments were promising, no stability analysis was provided for this earlier method, that would cope with errors generated by computer arithmetic or by inexact matrix-vector products. We show below that the new variant proposed in this paper are, by contrast, stable for fairly general classes of perturbations.

The paper is organized as follows. The range-space Krylov methods are derived for the symmetric and unsymmetric cases in Section 2. The effect of inexact matrix products on the convergence of these new algorithmic variants is then investigated in Section 3, and the concepts are numerically illustrated in Section 4. Some conclusions and perspectives are finally presented in Section 5.

2 Range-space Krylov methods

We start by considering the standard GMRES (Saad and Schultz, 1986) method for the system (1.1) for the matrix

$$A = \gamma I_n + K^T L. \quad (2.1)$$

Although this method is well-known, we briefly review its main features and concepts to establish a basis for later developments. The main idea of the method is to minimize the Euclidean norm of residual $(\gamma I_n + K^T L)s - b$ on the successive nested Krylov subspaces generated by the sequence

$$b, (\gamma I_n + K^T L)b, (\gamma I_n + K^T L)^2 b, (\gamma I_n + K^T L)^3 b, \dots \quad (2.2)$$

or, equivalently, by

$$b, (K^T L)b, (K^T L)^2 b, (K^T L)^3 b, \dots \quad (2.3)$$

This is achieved by using the Arnoldi process (see Saad, 1996, page 154, or Kelley, 1995, page 37) to generate nested orthonormal basis of these subspaces, i.e. a set of vectors $\{v_i\}_{i=1}^{k+1}$ with $v_1 = b/\|b\|$ (where $\|\cdot\|$ stands for the Euclidean norm) such that, after k steps,

$$(\gamma I_n + K^T L) V_k = V_{k+1} H_k, \quad (2.4)$$

where the columns of $V_k \stackrel{\text{def}}{=} [v_1, \dots, v_k]$ form an orthonormal basis of the k -th Krylov subspace

$$\mathcal{K}_k = \text{span}[b, \dots, (K^T L)^{k-1} b],$$

and where H_k is a $(k+1) \times k$ upper-Hessenberg matrix. The linear least-squares

$$\min_y \|H_k y - \beta_1 e_1\| \quad (2.5)$$

(where the symbol e_i denotes the i -th vector of the canonical basis and $\beta_1 = \|b\|$) is then explicitly solved in \mathbb{R}^k , yielding the GMRES⁽⁶⁾ algorithm on the following page.

In this statement of GMRES, one recognizes the construction of the Krylov sequence in Steps 1 and 2(a), its orthogonalization in Steps 2(b,d,e), the shift corresponding to the equivalence between (2.2) and (2.3) in Step 2(c) and the solution of the restricted linear

⁽⁶⁾Not to be confused with GMRES(m), a version which is restarted every m iterations.

Algorithm 2.1: $s = \text{GMRES}(K, L, b, \gamma)$

1. Define $\beta_1 = \|b\|$ and $v_1 = b/\beta_1$.
2. For $k = 1, \dots, m$,
 - (a) $w_k = K^T L v_k$
 - (b) for $i = 1, \dots, k$,
 - i. $H_{i,k} = v_i^T w_k$
 - ii. $w_k \leftarrow w_k - H_{i,k} v_i$
 - (c) $H_{k,k} \leftarrow H_{k,k} + \gamma$,
 - (d) $\beta_{k+1} = \|w_k\|$,
 - (e) $v_{k+1} = w_k / \beta_{k+1}$,
 - (f) $y_k = \arg \min_y \|H_k y - \beta_1 e_1\|$,
 - (g) if $\|H_k y_k - \beta_1 e_1\| < \epsilon_r$, break.
3. Return $s = V_k y_k$.

least-squares (2.5) in Step 2(f). The convergence test of Step 2(g) is there to detect early termination, but one knows that the residual $q_k = H_k y_k - \beta_1 e_1$ must decrease in norm at every iteration since $V_{k+1} q_k$ is nothing but the representation in \mathcal{K}_k of the projection of the right-hand side onto the orthogonal to \mathcal{K}_k . We refer the reader to Saad and Schultz (1986) for notation and convergence analysis.

If b does not belong to the range of A given by (2.1), there is no reason for q_k to converge to zero. Moreover it may happen in this case that “breakdown” occurs in the very unlikely circumstance where $A^k b$ turns out to be, for some k , an exact linear combination of $\{A^i b\}_{i=1}^{k-1}$, and the vector w_k may then be identically zero at the end of Step 2(b), thereby yielding $\beta_{k+1} = 0$ and making Step 2(e) undefined. The most obvious strategy to cope with this (mostly theoretical) situation is a simple restart of the process with a slightly perturbed right-hand side. Another interesting possibility is described in Reichel and Ye (2005), where an expanded Krylov space is built when a near-breakdown occurs, making it possible to continue the process.

The next step follows Gratton and Tshimanga (2009) and assumes (for now, see (2.10) and the discussion thereafter) that

$$b \in \text{range}(K^T). \quad (2.6)$$

We now observe that, in this case, the sequence (2.3) may be rewritten as

$$K^T d, \quad K^T (LK^T) d, \quad K^T (LK^T)^2 d, \quad K^T (LK^T)^3 d, \quad \dots \quad (2.7)$$

for some vector $d \in \mathbb{R}^m$, and thus deduce that the Krylov spaces in \mathbb{R}^m associated with this sequence are the images by K^T of other Krylov spaces generated now in \mathbb{R}^m by the sequence

$$d, \quad (LK^T) d, \quad (LK^T)^2 d, \quad (LK^T)^3 d, \quad \dots \quad (2.8)$$

We may therefore consider an Arnoldi process based on this sequence, leading now to the relation

$$LK^T \hat{V}_k \stackrel{\text{def}}{=} LK^T [\hat{v}_1, \dots, \hat{v}_k] = \hat{V}_{k+1} H_k \quad (2.9)$$

instead of (2.4), where we use the \hat{x} notation to denote a pre-image by K^T in \mathbb{R}^m of $x \in \text{range}(K^T) \subseteq \mathbb{R}^n$ (whose existence we show), i.e. some \hat{x} such that $x = K^T \hat{x}$. We now choose to use the (semi-)metric⁽⁷⁾ induced by KK^T , whose effect is to keep the modified

⁽⁷⁾We consider a semi-metric if K is rank deficient.

Arnoldi process mathematically equivalent to the standard one, modulo premultiplication by K^T . If we now rewrite the complete GMRES algorithm in (the much smaller) \mathbb{R}^m , we then obtain its variant RSGMR.

Algorithm 2.2: $s = \text{RSGMR}(K, L, d, \gamma)$

1. Define $p_1 = K^T d$, $\hat{z}_1 = K p_1$,
2. Set $\beta_1 = \sqrt{d^T \hat{z}_1}$, $\hat{v}_1 = d/\beta_1$, $\hat{z}_1 \leftarrow \hat{z}_1/\beta_1$ and $p_1 \leftarrow p_0/\beta_1$.
3. For $k = 1, \dots, m$,
 - (a) $\hat{w}_k = L p_k$
 - (b) for $i = 1, \dots, k$,
 - i. $H_{i,k} = \hat{z}_i^T \hat{w}_k$
 - ii. $\hat{w}_k \leftarrow \hat{w}_k - H_{i,k} \hat{v}_i$
 - (c) $H_{k,k} \leftarrow H_{k,k} + \gamma$,
 - (d) $p_{k+1} = K^T \hat{w}_k$, $\hat{z}_{k+1} = K p_k$, $\beta_{k+1} = H_{k+1,k} = \sqrt{\hat{z}_{k+1}^T \hat{w}_k}$,
 - (e) $\hat{v}_{k+1} \leftarrow \hat{w}_k/\beta_{k+1}$, $\hat{z}_{k+1} \leftarrow \hat{z}_k/\beta_{k+1}$, $p_{k+1} \leftarrow p_k/H_{k+1,k}$,
 - (f) $y_k = \arg \min_y \|H_k y - \beta_1 e_1\|$,
 - (g) if $\|H_k y_k - \beta_1 e_1\| < \epsilon_r$, break.
4. Return $s = K^T \hat{V}_k y_k$.

Observe that we had to compute the product of p_k by K , in order to evaluate the norms (in Steps 1 and 2(d)) and the inner products (in Step 2(b.i)) in the correct metric (i.e. in \mathbb{R}^n). Thus we have replaced computation and storage of a set n -dimensional vectors by that of two sets of m -dimensional ones for the price of an additional product by K at each iteration, with the exception of p_k which remains a “large” n -dimensional vector. RSGMR also requires the storage of the \hat{v}_i and the \hat{z}_i , but these are now of dimension m . Hence our comment in the introduction saying that work and storage depend essentially on m . Although the need for an additional product does not sound ideal, it may make the difference between an impractical method (where a collection of huge vectors is just too large for the computer at hand) and a more CPU-intensive but practical one. Observe finally that, since GMRES and RSGMR are mathematically equivalent, the sequence $\|q_k\|$ generated by RSGMR is identical to that generated by GMRES for $b = K^T d$ and therefore enjoys the same monotonicity property.

We have derived Algorithm RSGMR under the condition (2.6), which is, as we already mentioned, very commonly occurring in the context of the solution of the nonlinear problem (1.2). However, it is not difficult to avoid this assumption and allow for more generality by considering a right-hand side which is not the range of K^T . This is simply achieved by considering the extended system

$$(\gamma I + \bar{K}^T \bar{L})s = \bar{K}^T e_{m+1}, \quad \text{where } \bar{K} = \begin{bmatrix} K \\ b^T \end{bmatrix} \text{ and } \bar{L} = \begin{bmatrix} L \\ 0^T \end{bmatrix}, \quad (2.10)$$

since then

$$\bar{K}^T \bar{L} = K^T L \quad \text{and} \quad \bar{K}^T e_{m+1} = b. \quad (2.11)$$

These relations imply that the Krylov subspaces generated in \mathbb{R}^n when considering (2.10) are identical to those generated by RSGMR in \mathbb{R}^n , and therefore that the sequence of residuals norms produced by this reformulation is unaltered. Thus considering (2.10) does

not alter the variational properties of the original GMRES, as was already the case for RSGMR.

Because range-space methods are constrained to use the products by LK^T or KK^T , their numerical performance should not be expected to compare with that of “square root” methods (like the full-space LSQR) on severely ill-conditioned problems. While this is not a strong restriction for data assimilation applications (where the condition number of $I_n + \mathcal{H}^T \mathcal{H}$ rarely exceeds 10^4 , see Tshimanga, Gratton, Weaver and Sartenaer, 2008), this certainly reflects that substantially cheaper algorithms also come at a price. . . Fortunately, we show below that the accuracy obtained with RSGMR remains very acceptable in many cases⁽⁸⁾.

We conclude this section by a brief comparison of the computational and storage costs associated with full- and range-space GMRES. Table 2.1 reports the storage and computational costs⁽⁹⁾ at iteration k of these algorithms, where the computational cost is divided in operations internal to the algorithm and the products which have to be computed (very often using external software). In this table, the terms $k(k+3)/2$ in the storage costs correspond to storing the Hessenberg matrix H_k and y_k , while the first terms in the internal flops counts correspond to the cost of the relevant orthogonalization process. The symbol $[sol]$ represents the cost of solving the linear least-squares (2.5) in the k -dimensional Krylov subspace. We also give in Table 2.2 the initialization and termination computational costs, assuming that termination occurs at iteration k .

	GMRES	RSGMR
storage	$n(k+1) + k(k+3)/2$	$n + (2m+1)k + k(k+3)/2$
internal flops	$4nk + 3n + [sol]$	$4mk + 7m + [sol]$
products by	K^T, L	K^T, K, L

Table 2.1: Comparative storage, floating point operations and matrix vector products at iteration k for full- and range-space GMRES

	GMRES	RSGMR
initialization	$3n$	$2n + 2m + \text{prod}(K)$
termination (k)	$2nk$	$2(m+1)k + \text{prod}(K^T)$

Table 2.2: Comparative initialization and termination computational costs for full- and range-space GMRES

The comparative advantages and drawbacks of the range-space method appear clearly in these tables: the range-space variant is preferable when $m \ll n$ and when the cost of an additional product by K per iteration is not prohibitive compared to that of reduced storage.

3 Convergence with inexact products

After deriving RSGMR, we now propose an analysis of the behaviour of Krylov algorithms in the case where the products by K^T and L are performed inexactly. The main reason to consider inexact matrix products is clearly to allow them in an inexact but cheaper form. This is especially crucial in the context of data assimilation as discussed in Section 1 where the products constitute by far the most costly part of the computation. In particular, inexact products allow us to consider, in this context, inexact solves with B^{-1} or R^{-1} as well as the use of less expensive, degraded versions of the operators \mathcal{H} and \mathcal{H}^T . Similar

⁽⁸⁾Extensive numerical experiments not reported here show that the best achievable accuracy for RSGMR is comparable to that obtained with normal-equations approaches.

⁽⁹⁾We ignore the constants in this evaluation.

considerations, albeit in different contexts, were also used for motivating the analysis proposed by van den Eshof and Sleijpen (2004) and by Simoncini and Szyld (2003) for the full-space versions of iterative Krylov solvers.

In the following analysis, we assume that the first products (corresponding to the scaling of d) are already performed inexactly, that is

$$\hat{v}_1 = d/\beta_1 = d/\sqrt{d^T(K + E_{K,0})(K^T + E_{K^T,0})d} \quad (3.1)$$

for some error matrices $E_{K^T,0}$ and $E_{K,0}$. Assume also that each subsequent product by K^T , K or L is inexact in the sense that, at iteration i ,

$$L_i = L + E_{L,i}, \quad K_i^T = K^T + E_{K^T,i} \quad \text{and} \quad K_i = K + E_{K,i},$$

for some errors $E_{L,i}$, $E_{K^T,i}$ and $E_{K,i}$. If we terminate at iteration k , the solution $s_k = s$ is finally computed using an inexact product with K^T in the formula

$$s_k = K_*^T \hat{V}_k y_k \quad \text{where} \quad K_* \stackrel{\text{def}}{=} K + E_{K^T,*}^T \quad (3.2)$$

for some error $E_{K^T,*}$.

We now propose two different models for describing the inaccuracy in the matrix-vector products. In the first model, which we call the *backward-error model*, we assume that

$$\|E_{K,i}\| \leq \tau_{K,i}\|K\|, \quad \text{and} \quad \|E_{K^T,i}\| \leq \tau_{K^T,i}\|K\| \quad \text{for } i = 0, \dots, k, \quad (3.3)$$

$$\|E_{L,i}\| \leq \tau_{L,i}\|L\| \quad \text{for } i = 1, \dots, k, \quad \text{and} \quad \|E_{K^T,*}\| \leq \tau_*\|K\|$$

for some tolerances $\tau_{K,i}$, $\tau_{K^T,i}$, $\tau_{L,i}$ and τ_* belonging to the interval $[0, 1)$.

The second error model for inexact products, called the *forward-error model* is stronger and replaces bounds on the errors on the operators by bounds on the errors on the vector resulting from the application of the operator. In this model, we replace the above bounds by

$$\|E_{K,i} u_n\| \leq \tau_{K,i}\|K u_n\|, \quad \text{and} \quad \|E_{K^T,i} u_m\| \leq \tau_{K^T,i}\|K u_m\| \quad \text{for } i = 0, \dots, k, \quad (3.4)$$

$$\|E_{L,i} u_n\| \leq \tau_{L,i}\|L u_n\| \quad \text{for } i = 1, \dots, k, \quad \text{and} \quad \|E_{K^T,*} u_m\| \leq \tau_*\|K u_m\|$$

where u_n and u_m are vectors of dimension n and m , respectively, to which the operators K or L (for u_n) or K^T (for u_m) are applied.

Which error model is preferable is unclear in general and might depend on context. While the backward-error approach is more widespread in the literature (it used by van den Eshof and Sleijpen, 2004, and Simoncini and Szyld, 2003), the forward-error approach may be judged more realistic in situations where monitoring the output of a complex process for a specific input is feasible, but impossible or too expensive for all possible inputs.

Our aim is then to bound $\|r_k\|$, the norm of the true residual at iteration k , where

$$r_k = (\gamma I_n + K^T L)s_k - K^T d,$$

if possible by quantities which can be obtained or estimated in the course of the computation. We first conduct our analysis under the assumption that no breakdown occurs, that is

$$\hat{w}_i \neq 0 \quad \text{for } i = 1, \dots, k, \quad (3.5)$$

(where \hat{w}_i is considered at the end of the normalization, that is after Step 2(b)), but we will comment on the situation where this condition fails at the end of the section.

We start by analyzing the perturbed Arnoldi iteration (see Golub and Van Loan, 1989, p. 499) and prove a useful bound on the residual norm, irrespective of the error model considered.

Lemma 3.1 *Suppose that no breakdown occurs. Then one has that*

$$\|r_k\| \leq \|Q_k\| \|H_k y_k - \beta_1 e_1\| + \gamma \|E_{K^T, *} \hat{V}_k y_k\| + \|K\| \sum_{i=1}^k |[y_k]_i| \|(LK_*^T - L_i K_i^T) \hat{v}_i\| \quad (3.6)$$

where $Q_k \stackrel{\text{def}}{=} K^T \hat{V}_{k+1}$.

Proof. By constuction, RSGMR ensures that the Arnoldi relation (2.9) holds with perturbed matrices, that is

$$[(\gamma I_n + L_1 K_1^T) \hat{v}_1, \dots, (\gamma I_n + L_k K_k^T) \hat{v}_k] = \hat{V}_{k+1} H_k. \quad (3.7)$$

Observe now that, because of (3.1),

$$K^T d = \beta_1 K^T \hat{v}_1 = \beta_1 K^T \hat{V}_{k+1} e_1.$$

Subtracting this quantity from both sides of (3.7) premultiplied by K^T , we obtain that

$$K^T [(\gamma I_n + L_1 K_1^T) \hat{v}_1, \dots, (\gamma I_n + L_k K_k^T) \hat{v}_k] y_k - K^T d = K^T \hat{V}_{k+1} (H_k y_k - \beta_1 e_1).$$

Remembering now that s_k is given by (3.2), we see that

$$\begin{aligned} \|r_k\| &= \|K^T L K_*^T \hat{V}_k y_k - K^T d + \gamma K_*^T \hat{V}_k y_k\| \\ &= \|K^T L K_*^T \hat{V}_k y_k + K^T \hat{V}_{k+1} (H_k y_k - \beta_1 e_1) \\ &\quad - K^T [(\gamma I_n + L_1 K_1^T) \hat{v}_1, \dots, (\gamma I_n + L_k K_k^T) \hat{v}_k] y_k + \gamma K_*^T \hat{V}_k y_k\| \\ &= \|K^T L K_*^T \hat{V}_k y_k + Q_k (H_k y_k - \beta_1 e_1) + \gamma (K_*^T - K^T) \hat{V}_k y_k \\ &\quad - K^T [L_1 K_1^T \hat{v}_1, \dots, L_k K_k^T \hat{v}_k] y_k\| \\ &\leq \|Q_k\| \|H_k y_k - \beta_1 e_1\| + \gamma \|E_{K^T, *} \hat{V}_k y_k\| \\ &\quad + \|K\| \|L K_*^T \hat{V}_k y_k - [L_1 K_1^T \hat{v}_1, \dots, L_k K_k^T \hat{v}_k] y_k\| \end{aligned} \quad (3.8)$$

But we have, using the triangle inequality, that

$$\|L K_*^T \hat{V}_k y_k - [L_1 K_1^T \hat{v}_1, \dots, L_k K_k^T \hat{v}_k] y_k\| \leq \sum_{i=1}^k |[y_k]_i| \|L K_*^T \hat{v}_i - L_i K_i^T \hat{v}_i\|$$

and (3.6) follows. \square

After this general result, we now focus on the backward-error model.

Lemma 3.2 *Suppose that the backward-error model holds and that no breakdown occurs, one has that*

$$\|r_k\| \leq \|Q_k\| \|H_k y_k - \beta_1 e_1\| + \|K\| \pi_k \left[\tau_{K^T, *} \gamma \sqrt{k} \|y_k\| + 4 G^2 \sum_{i=1}^k |[y_k]_i| \tau_i \right], \quad (3.9)$$

where $\pi_k \stackrel{\text{def}}{=} \max_{i=1, \dots, k} \|\hat{v}_i\|$, $G \stackrel{\text{def}}{=} \max[\|K\|, \|L\|]$ and $\tau_i \stackrel{\text{def}}{=} \max[\tau_*, \tau_{K^T, i}, \tau_{K, i}, \tau_{L, i}]$.

Proof. We first obtain, using the triangle inequality and the definition of π_k , that

$$\begin{aligned} \sum_{i=1}^k |[y_k]_i| \|(L K_*^T - L_i K_i^T) \hat{v}_i\| &\leq \sum_{i=1}^k |[y_k]_i| \|L K_*^T - L_i K_i^T\| \|\hat{v}_i\| \\ &\leq \pi_k \sum_{i=1}^k |[y_k]_i| \|L K_*^T - L_i K_i^T\|. \end{aligned}$$

Again using the triangle inequality and (3.3), we now deduce that

$$\begin{aligned} \|LK_*^T - L_i K_i^T\| &= \|LE_{K^T,*} - LE_{K^T,i} - E_{L,i}K^T - E_{L,i}E_{K^T,i}\| \\ &\leq 3\tau_i G^2 + \tau_i^2 G^2 \\ &\leq 4\tau_i G^2, \end{aligned} \quad (3.10)$$

where we used the bound $\tau_i \leq 1$ to derive the last inequality, and therefore that

$$\sum_{i=1}^k |[y_k]_i| \|(LK_*^T - L_i K_i^T)\hat{v}_i\| \leq 4\pi_k G^2 \sum_{i=1}^k |[y_k]_i| \tau_i.$$

Substituting this bound in (3.6) and using the inequality

$$\|E_{K^T,*}\hat{V}_k y_k\| \leq \tau_* \|K\| \pi_k \sqrt{k} \|y_k\| \leq \tau_i \|K\| \pi_k \sqrt{k} \|y_k\|$$

then gives (3.9). \square

Observe that the proof of this lemma does not use $\tau_{k,i}$ and that the error on the products by K does not explicitly appear in the bound (3.9), but is present implicitly as the quantities $\|Q_k\|$ and π_k crucially depend on the metric KK^T and therefore, in our case, on these errors $E_{K,i}$ ($i = 1, \dots, k$). We now bound these quantities, provided the error remains sufficiently small compared to the condition number of K defined by

$$\kappa(K) \stackrel{\text{def}}{=} \frac{\sigma_{\max}(K)}{\sigma_{\min}^0(K)}$$

where $\sigma_{\max}(K)$ is the largest singular value of K and $\sigma_{\min}^0(K)$ is the smallest of the strictly positive ones (e.g., Björck, 1996, page 28, ...).

Lemma 3.3 *Suppose that the backward-error model holds, that no breakdown occurs and that is sufficiently small to ensure that, for all i ,*

$$\tau_i \kappa(K) < \frac{1}{6}. \quad (3.11)$$

Then

$$\|Q_k\| \leq \sqrt{2(k+1)} \quad \text{and} \quad \pi_k \leq \frac{\sqrt{2}}{\sigma_{\min}^0(K)}. \quad (3.12)$$

Proof. From the Cauchy-Schwarz inequality, the fact that the vectors \hat{v}_i are orthonormal in the KK^T induced inner product by construction and (3.3), we first verify that

$$\frac{|\hat{v}_i^T KE_{K^T,i}\hat{v}_i|}{\hat{v}_i^T KK^T \hat{v}_i} \leq \frac{\|E_{K^T,i}\hat{v}_i\|}{\|K^T \hat{v}_i\|} \leq \frac{\tau_{K^T,i} \|K\| \|\hat{v}_i\|}{\sigma_{\min}^0(K) \|\hat{v}_i\|} \leq \tau_i \kappa(K). \quad (3.13)$$

Similarly,

$$\frac{|\hat{v}_i^T E_{K,i}K^T \hat{v}_i|}{\hat{v}_i^T KK^T \hat{v}_i} \leq \tau_i \kappa(K) \quad \text{and} \quad \frac{|\hat{v}_i^T E_{K,i}E_{K^T,i}\hat{v}_i|}{\hat{v}_i^T KK^T \hat{v}_i} \leq \tau_i^2 \kappa(K)^2. \quad (3.14)$$

Thus we deduce from (3.11), (3.13) and (3.14) that

$$\left| \frac{\hat{v}_i^T (KE_{K^T,i} + E_{K,i}K^T + E_{K,i}E_{K^T,i})\hat{v}_i}{\hat{v}_i^T KK^T \hat{v}_i} \right| \leq 3\tau_i \kappa(K) \leq \frac{1}{2}. \quad (3.15)$$

As a consequence, since clearly

$$\hat{v}_i^T (K + E_{K,i})(K^T + E_{K^T,i})\hat{v}_i = \left[1 + \frac{\hat{v}_i^T (KE_{K^T,i} + E_{K,i}K^T + E_{K,i}E_{K^T,i})\hat{v}_i}{\hat{v}_i^T KK^T \hat{v}_i} \right] \hat{v}_i^T KK^T \hat{v}_i,$$

we thus obtain that

$$\frac{1}{2}\|K^T \hat{v}_i\|^2 \leq \hat{v}_i^T (K + E_{K,i})(K^T + E_{K^T,i}) \hat{v}_i \leq \frac{3}{2}\|K^T \hat{v}_i\|^2 \quad (3.16)$$

and the normalization (with respect to the inexact metric $K_i K_i^T$) performed by the algorithm to ensure that

$$\hat{v}_i^T (K + E_{K,i})(K^T + E_{K^T,i}) \hat{v}_i = 1 \quad (3.17)$$

is legal because of (3.5) and because \hat{v}_i belongs to the range of K^T by construction. Inserting (3.17) in (3.16), we then deduce that

$$\frac{2}{3} < \hat{v}_i^T K K^T \hat{v}_i < 2. \quad (3.18)$$

and (3.12) follows. \square

While the bound on $\|Q_k\|$ given by (3.12) is formally correct and conceptually tight, it is often very pessimistic in practice, as it does not take the typical random nature of the error into account. In particular, it is not unusual for the residual error to be independent of the factor $\sqrt{k+1}$ in the bound expression.

We may then combine the two above lemmas and obtain a final set of bounds on $\|r_k\|$ for the case where the backward-error model is considered.

Theorem 3.4 *Suppose that the backward-error model holds, that no breakdown occurs and that $\tau_{\max} \stackrel{\text{def}}{=} \max[\tau_1, \dots, \tau_k]$ is small enough to ensure (3.11). Then*

$$\begin{aligned} \|r_k\| &\leq \sqrt{2(k+1)} \|H_k y_k - \beta_1 e_1\| + \|K\| \pi_k \left[\tau_* \gamma \sqrt{k} \|y_k\| + 4 G^2 \sum_{i=1}^k |[y_k]_i| \tau_i \right] \\ &\leq \sqrt{2(k+1)} \|H_k y_k - \beta_1 e_1\| + \tau_{\max} \sqrt{k} \|K\| \pi_k (\gamma + 4 G^2) \|y_k\| \\ &\leq \sqrt{2(k+1)} [\|H_k y_k - \beta_1 e_1\| + \tau_{\max} \kappa(K) (\gamma + 4 G^2) \|y_k\|]. \end{aligned} \quad (3.19)$$

The first and sharpest of these three bounds allows the consideration of levels of inexactness in the products that vary from iteration to iteration, in the spirit of Simoncini and Szyld (2003) and van den Eshof and Sleijpen (2004). It only involves computable quantities, as $\|H_k y_k - \beta_1 e_1\|$ (the norm of the residual in \mathcal{K}_k), $\|y_k\|$ and π_k can be recurred within the RSGMR iterations in $O(\max[m, k^2])$ operations⁽¹⁰⁾. The second bound is most interesting in the case where a constant bound on inexactness is selected. The third bound is similar, but does not requires the computation of π_k , at the cost of a (experimentally often severe) overestimation.

We observe the specific role of τ_* in the error bound, as it appears multiplied by $\|y_k\|$ rather than by one of the $|[y_k]_i|$, a phenomenon specific to the range-space setting. The final product by K^T to produce s in Step 4 should therefore be computed with a potentially higher accuracy than any of the preceding products if the overall error bound is to be preserved. It is also interesting to note that there is little to be gained by controlling the errors of the products of K , K^T and L differently, as the error-bound analysis show their effects to be intertwined (see e.g. (3.15)). This is also born out in our numerical experiments (see Section 4).

Deriving similar bounds for the case where the forward-error model holds can be achieved along the same lines. The first salient difference with the analysis presented for the backward-error model is that we may now bound $\|(LK_*^T - L_i K_i^T) \hat{v}_i\|$ in (3.6) directly, rather than $\|LK_*^T - L_i K_i^T\|$ and $\|\hat{v}_i\|$ separately (as in Lemma 3.2), which avoids the introduction of the square power G and of π_k in the resulting bounds (3.10) and (3.9).

⁽¹⁰⁾ Assuming the use of Givens rotations to compute y_k .

The second difference is that the denominator in the second term of (3.13) no longer needs to be bounded below and simplifies with the numerator using the forward-error bound (3.4), thereby removing the dependence of $\|Q_k\|$ from $\kappa(K)$. Taking these differences into account and using the inequality $\|K\| \leq G$ in (3.6) then yields the following result.

Theorem 3.5 *Suppose that the forward-error model holds, that no breakdown occurs and that $\tau_{\max} < \frac{1}{6}$. Then one has that*

$$\begin{aligned} \|r_k\| &\leq \sqrt{2(k+1)} \|H_k y_k - \beta_1 e_1\| + \sqrt{2} \left[\tau_* \gamma \sqrt{k} \|y_k\| + 4G^2 \sum_{i=1}^k |[y_k]_i| \tau_i \right] \\ &\leq \sqrt{2(k+1)} \left[\|H_k y_k - \beta_1 e_1\| + \tau_{\max} (\gamma + 4G^2) \|y_k\| \right]. \end{aligned} \quad (3.20)$$

A comparison of Theorems 3.4 and 3.5 shows that the second is considerably stronger, in that the conditioning of K or the factor $\|K\|\pi_k$ no longer appear. As a consequence, possibly significantly more inexact products are possible in the forward-error model than in the backward-error one, with the same final accuracy requirement.

We observe that breakdown may obviously occur in the inexact case just as with exact products, and can be solved using the same strategies. In the context of inexact products, a simpler method also consists in recomputing the last matrix-vector product with a marginally different accuracy threshold, which then corresponds to a slightly different error in the inexact product. This suppresses the breakdown (with probability one) and allows the algorithm to proceed.

We finally note that, because it uses exact arithmetic, our analysis applies without modification to mathematically equivalent variants of GMRES, including classical or iterated Gram-Schmidt procedures, as well as Householder variants.

4 Numerical illustrations

We now illustrate some of the concepts and motivation numerically and show the effect of using different strategies for allowing inexactness in the products in applying RSGMR to the system (1.1). The simplest technique is to decide of a product accuracy threshold in view of the desired final accuracy on the normalized true system residual, the latter being given (assuming termination at iteration k) by

$$\frac{\|r_k\|}{\|A\| \|s_*\|}, \quad (4.1)$$

where s_* is again the true system's solution. In the results presented below, we have chosen this threshold τ according to the formula

$$\tau_{BEM} = \frac{40\epsilon}{\sqrt{2(m+1)} \kappa(K)} \quad \text{and} \quad \tau_{FEM} = \epsilon,$$

where BEM and FEM refer to the backward-error model and the forward-error model, respectively. The formula for τ_{BEM} is inspired by the bound (3.19) (using the empirical observations that $\gamma + 4G^2 \approx 4G^2$ and $\|y_k\| = O(\|s_*\|)$) and a factor 10 to counteract the looseness of the bound). That for τ_{FEM} is directly derived from (3.20). The RSGMR algorithm using these two strategies and associated error models are applied to an example with $\gamma = 1$, $n = 1000$ and $m = 100$ whose matrices K and L are chosen randomly with m nonzero singular values whose logarithms are equally spaced between 0.1 and 0.3. The requested accuracy is chosen as $\epsilon = 10^{-5}$. The results are shown in Figure 4.1, in which the true normalized residual norms (4.1) are represented by a solid line, normalized Krylov residual norms

$$\frac{\|q_k\|}{\|H_k\| \|y_k\|}$$

by a dashed line, final accuracy requirement by an horizontal dotted line and accuracy threshold τ at iteration k by a dashed-dotted line (the same graphical convention is used in all subsequent figures). The left graph shows the effect of applying the backward-error model (with $\tau_k = \tau_{BEM}$) and the left graph that of applying the forward-error model (with $\tau_k = \tau_{FEM}$).

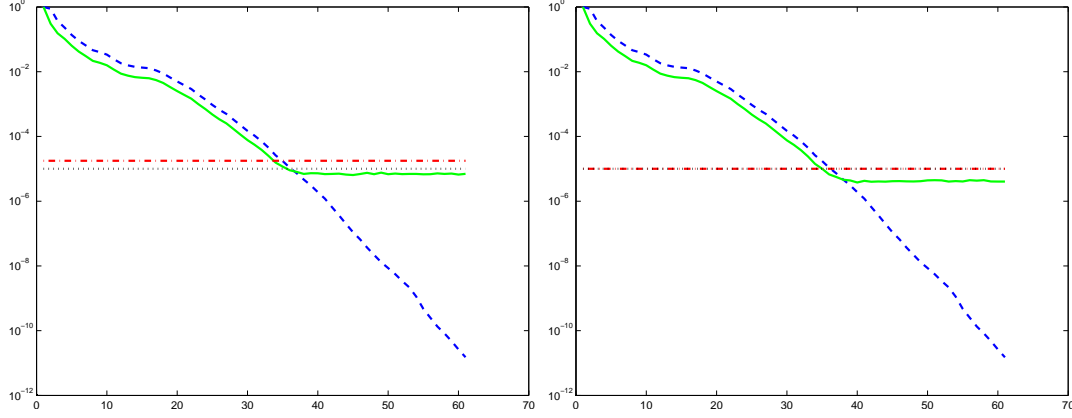


Figure 4.1: Accuracy threshold, normalized true and Krylov residual norms as a function of k when using the backward-error (left) and forward-error (right) accuracy models for the products by K , K^T and L ($\sigma_{\min}(A) \approx 2 \times 10^{-2}$, $\sigma_{\max}(A) \approx 4$). See the text for the description of each curve's meaning.

The differences between the two error strategies are small in this case. Being also interested in variable accuracy thresholds strategies, we also applied a technique recommended by Simoncini and Szyld (2003) (and adapted to (4.1)): we ran the same example with the choice

$$\tau_i = \tau_{SS} = \frac{\sigma_{\min}(A)}{m} \frac{\epsilon \|s_*\|}{\|q_i\|} \quad \text{and} \quad \tau_* = \tau_{BEM} \quad (4.2)$$

within the backward-error accuracy model, which is consistent with assumptions used in this reference. The result of this experiment is shown in Figure 4.2.

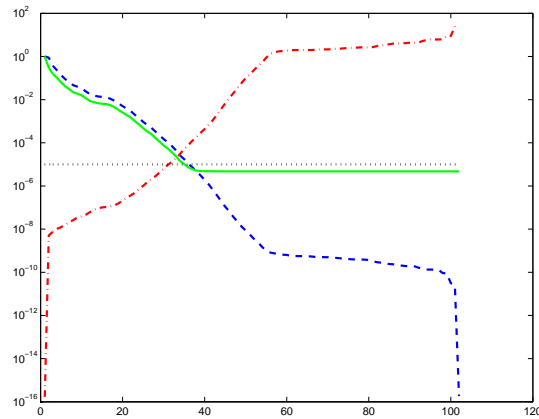


Figure 4.2: Accuracy threshold, normalized true and Krylov residual norms as a function of k when using the variable backward accuracy requirements described by (4.2) for the products by K , K^T and L ($\sigma_{\min}(A) \approx 2 \times 10^{-2}$, $\sigma_{\max}(A) \approx 4$). See the text for the description of each curve's meaning.

We see in this case a profile of the accuracy threshold similar to those reported by Simoncini and Szyld, but note that the variable threshold is significantly below the thresholds illus-

trated in Figure 4.1 practically up to the point where the maximal accuracy has been reached for the true residual.

We now apply the same algorithm on the same example modified so that the nonzero singular values of K and L have their logarithm equally spaced between 1 and 3 (instead of 0.1 and 0.3), significantly affecting both $\kappa(K)$ and $\|K\|$. The results are shown in Figure 4.3.

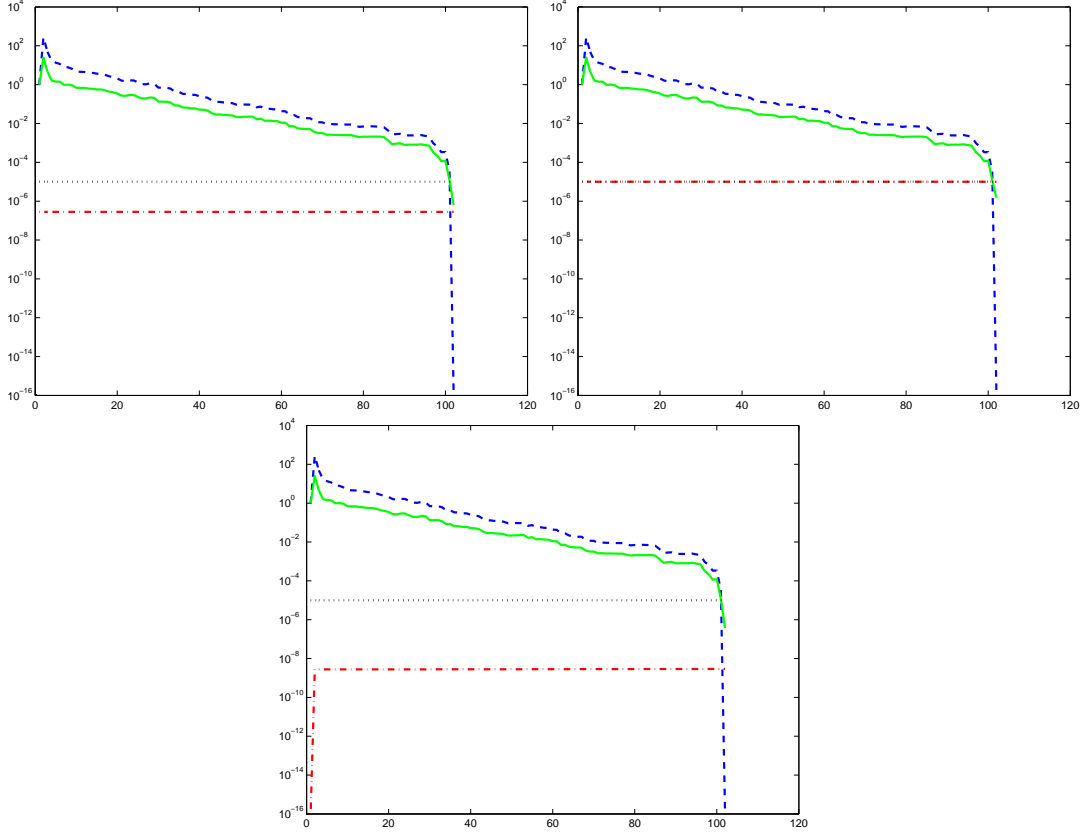


Figure 4.3: Accuracy threshold, normalized true and Krylov residual norms as a function of k when using the backward-error (top-left) and forward-error (top-right) accuracy models and the variable model (bottom) for the products by K , K^T and L ($\sigma_{\min}(A) \approx 2 \times 10^{-4}$, $\sigma_{\max}(A) \approx 5 \times 10^5$). See the text for the description of each curve's meaning.

As expected the accuracy threshold for the backward-error model decreases (to compensate the change in conditioning and $\|K\|$), but the forward-error-model threshold remains unchanged. Unfortunately, the variable strategy now picks up a too severe accuracy threshold because of the smaller $\sigma_{\min}(A)$ and only allows a small error on the products, thereby illustrating the difficulties of designing a robust and efficient variable accuracy scheme.

We finally conclude our numerical illustrations by verifying the claim made above that manipulating the accuracy thresholds on K and L differently does not affect the algorithm much. To this aim, we return to our second test case (for which the behaviour of RSGMR is shown in Figure 4.1), and run the algorithm first allowing inexact products with L only, and then with K and K^T only. Figure 4.4 presents the results of these two runs (using the backward-error model), and one checks that they do not differ significantly from the left graph of Figure 4.1 (except maybe a marginally lower final true residual when the product with K^T is exact, probably resulting from a better accuracy in recovering the final s).

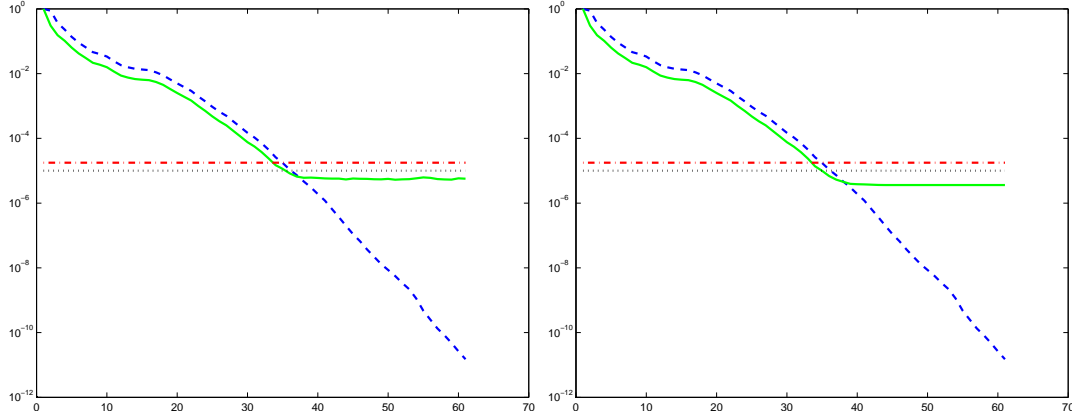


Figure 4.4: Impact of different accuracy thresholds for the products by K and L : no error on products by L on left and no error on products by K and K^T on the right. See the text for the description of each curve's meaning.

5 Conclusions and perspectives

Motivated by applications in inverse problems and, more specifically, by data assimilation for oceanography and wheather forecasting, we have introduced a range-space variant of GMRES. This variant is characterized by significantly lower storage requirements and inner computational costs than its full-space version, at the cost of an additional matrix vector product per iteration.

With the aim of reducing the computational burden further, we have also considered how strategies involving inexact matrix-vector products can be applied in the new variant, and have distinguished two distinct models to describe the inexactness allowed in these products. The first of this model considers the backward error on the matrix itself, while the second only controls the forward error on the vector resulting from the product. Formal error bounds on the true system residual are derived in both cases, indicating that the second strategy might allow looser accuracy requirements than the first.

We have finally provided numerical illustrations confirming these findings and also indicating that strategies in which the accuracy of the matrix-vector product varies in course of the iterations may be difficult to design while the simpler technique of using constant accuracy thresholds appears to work well. We may therefore conclude that range-space methods are remarkably efficient when applicable.

Although not covered explicitly in this paper, a range-space version of FOM (Full Orthogonalization Method) can be obtained from the derivation presented above by choosing $L = K^T$ and replacing the least-squares subproblem at Step 3(f) of RSGMR by the computation of the solution y_k of the system

$$H_k^\square y_k = \beta_1 e_1$$

where $e_1 \in \mathbb{R}^k$ and H_k^\square is the leading $k \times k$ submatrix of H_k . Interestingly, this variant has the further advantage of requiring exactly as many matrix-vector products per iteration as the original full-space FOM algorithm. Range-space versions of other Krylov methods such as MINRES, CG and LSQR may be derived along the same lines. For details of these derivations, we refer the reader to Gratton, Toint and Tshimanga (2009), where numerical evidence is also presented indicating that, if matrix-vector products are computed to machine accuracy, short recurrence methods such as range-space variants of CG or MINRES may be adequate, while full reorthogonalization techniques like range-space variants of FOM (Full Orthogonalisation Method) and GMRES are preferable if these products are computed more inexactly.

The authors are well aware that a complete numerical evaluation of the new range-space variants is necessary and that a number of issues raised in this paper merit further development. These include, in particular, the design of an efficient and robust variable accuracy scheme, clear stopping rules based on the formal error estimates, additional short recurrences techniques beyond CG and MINRES, and several other implementations issues such as the selection of the most suitable GMRES formulation amongst Gram-Schmidt and Householder variants. The extension of the forward-error model to derive tighter residual bounds for full-space methods may also be of interest. These topics are the object of ongoing research. A numerical comparison with restarted variants of Krylov methods may also be worthwhile to verify the advantage of full-orthogonalization techniques in the presence of inexact products.

It is of course especially worthwhile to apply the new range-space algorithms in the context of the motivating inverse problems in data assimilation. This longer-term work involves adapting new software libraries to the complex environment of weather forecasting systems and is currently ongoing. In this framework, it is also necessary to specialize the range-space Krylov variants further to gracefully handle the correlation matrices B and R . Inexact products are especially appealing in this case, where iterative solutions of linear systems involving these matrices may be truncated, or variable-fidelity techniques (like multigrid or simplified physics) exploited to alter the accuracy of the underlying models.

Acknowledgements. This research presented in this paper was conducted with the support of the “Assimilation de Données pour la Terre, l’Atmosphère et l’Océan (ADTAO)” project, funded by the Fondation “Sciences et Technologies pour l’Aéronautique et l’Espace (STAE)”, Toulouse, France, within the “Réseau Thématique de Recherche Avancée (RTRA)”.

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